

# ***A join point for loops in AspectJ***

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## *What we would like to do*

- Writing aspects that represent the concern:
  - “parallelise all the loops iterating from 0 to the length of an array of int using MPI”,
  - or “parallelise all the loops iterating over a Collection using Java Threads”.
- Write (aspect) code that does not invade the readability of the numerical code.

## ***Previously, on loops and AspectJ...***

- “*Using AspectJ to Separate Concerns In Parallel Scientific Java Code*” (AOSD 2004)
- Parallelisation of loops using aspects:
  - by making the iteration space visible as parameters to the methods
  - by turning loops into self-contained objects (loop body and boundaries )
- Both require refactoring the base code

# *Presentation Outline*

- Join point model:
  - Part 1: Shadows (static part),
  - Part 2: Context exposure (dynamic part),
- Loop selection,
- Implementation using `abc`,
- Dealing with exceptions,
- Related topics.

## *Join Points*

- A **join point** is “*a point in the dynamic call graph of a running program*”.
- A join point **shadow** is its location in the text of the program.
- Ability to weave code *before, after* and/or *around*.
- Ability to access **execution context**.

# ***JP Part 1: Shadows (static)***

- Analysis of the control flow graph
- Finding natural and combined loops
- Classification of loops according to their weaving and analysis capabilities:
  - General loops
  - Loops with unique successor
  - Loops with unique exit node

# ***Control-flow graph, dominators and natural loops (I)***

- A node is a **basic block** (only entry via its head and only exit via its tail).
- Node  $d$  **dominates** node  $n$  if every path from the beginning to  $n$  goes through  $d$ .
- A **back edge** ( $a \rightarrow b$ ) is an edge whose head ( $b$ ) dominates its tail ( $a$ ).
- Given a back edge  $n \rightarrow d$ , the natural loop is  $d$  plus the set of nodes that can reach  $n$  without going through  $d$ .

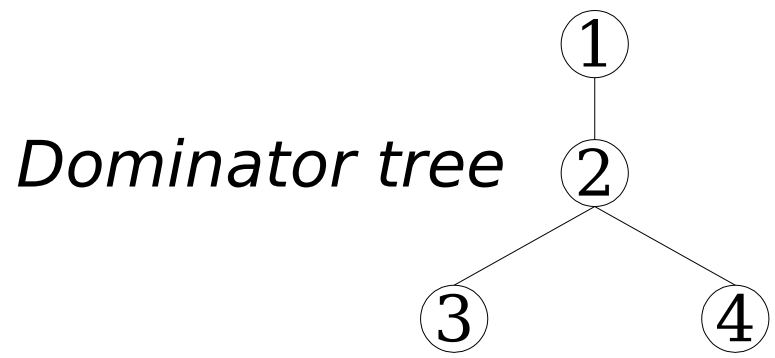
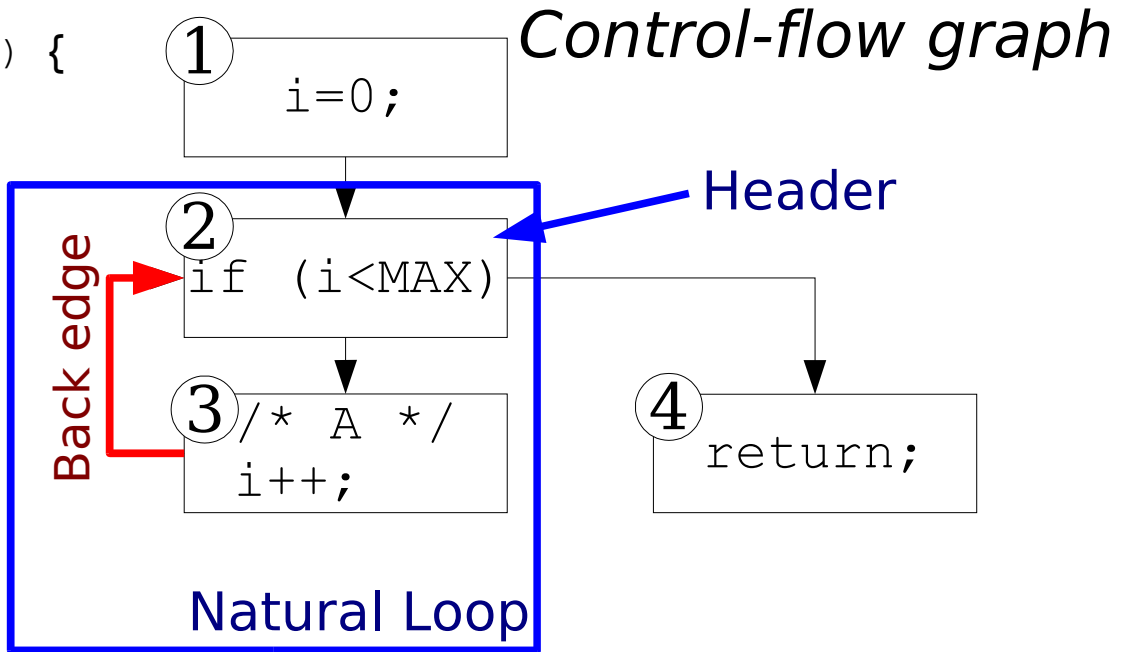
# Control-flow graph, dominators and natural loops (II)

```

for (int i = 0 ; i<MAX ; i ++ ) {
    /* A */
}

int j = 0 ;
int STRIDE = 1 ;
for (; j < MAX ; j+=STRIDE) {
    /* A */
}

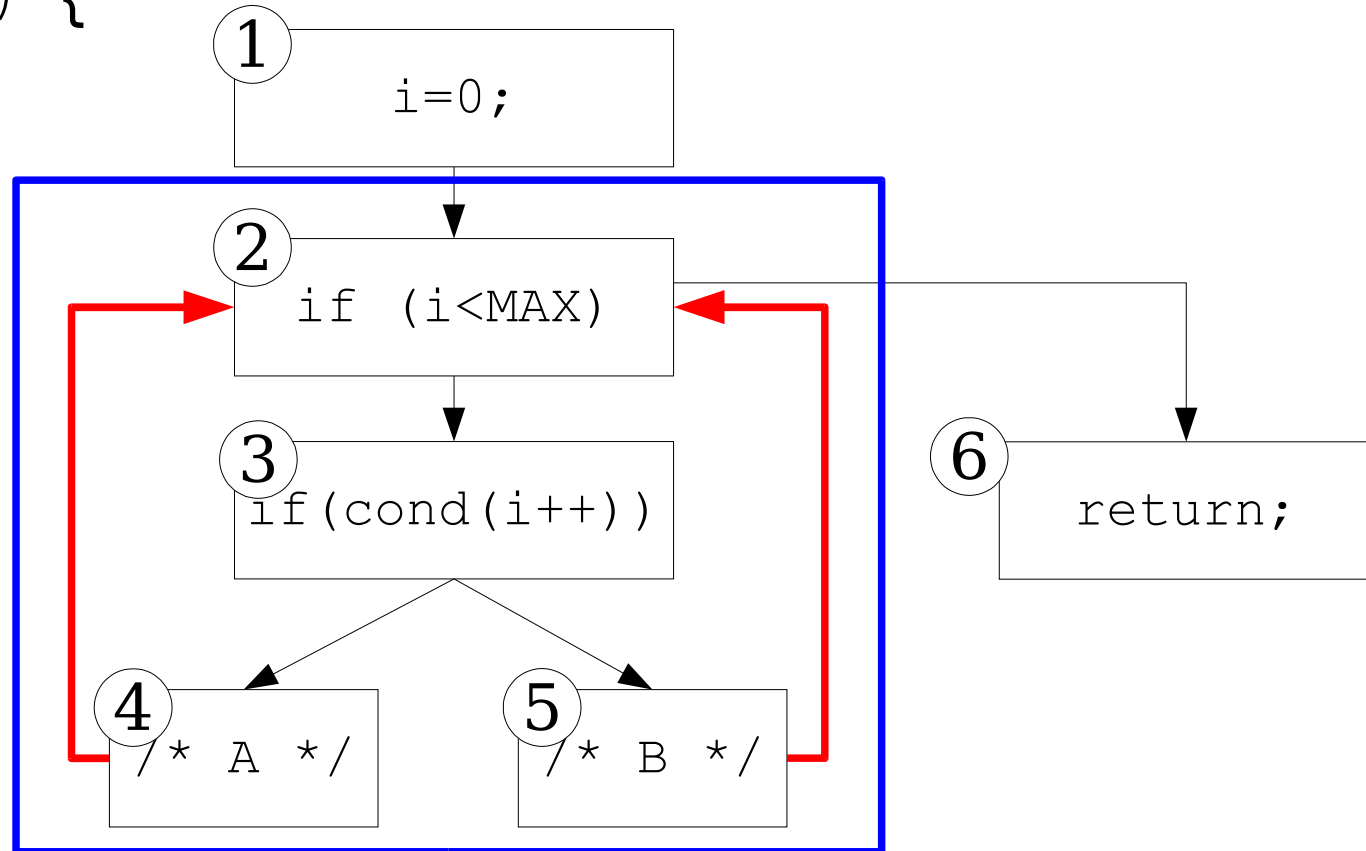
int k = 0 ;
while (k < MAX) {
    /* A */
    k ++ ;
}
    
```





# Combined loops

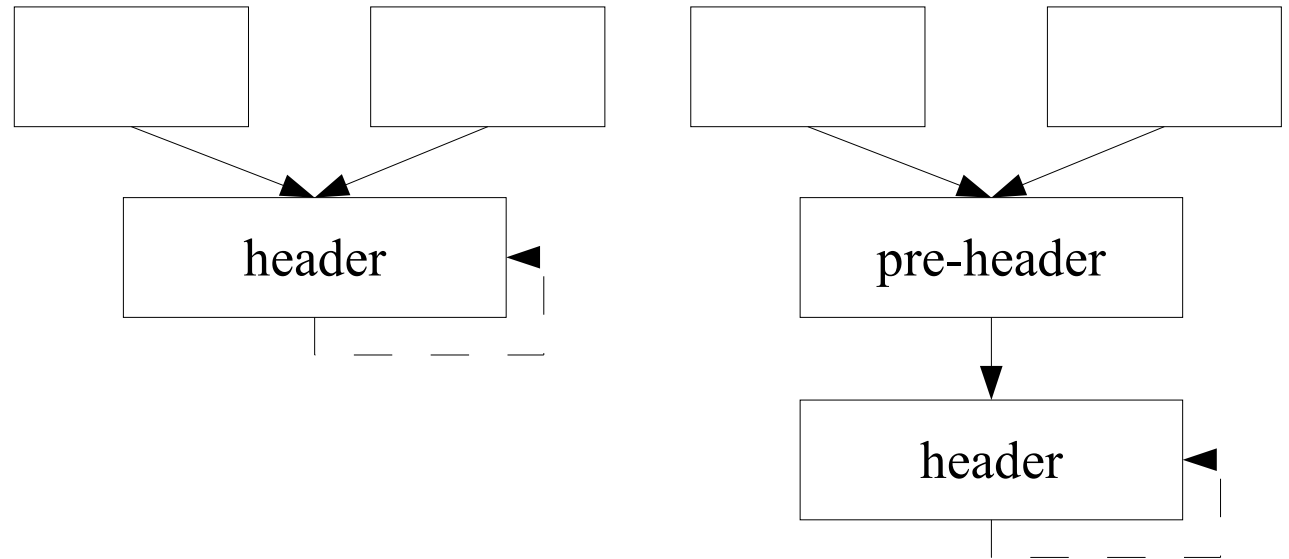
```
int i = 0 ;
while (i < MAX) {
    if (cond(i++)) {
        /* A */
    } else {
        /* B */
    }
}
```



1 combined loop  
with 2 back edges

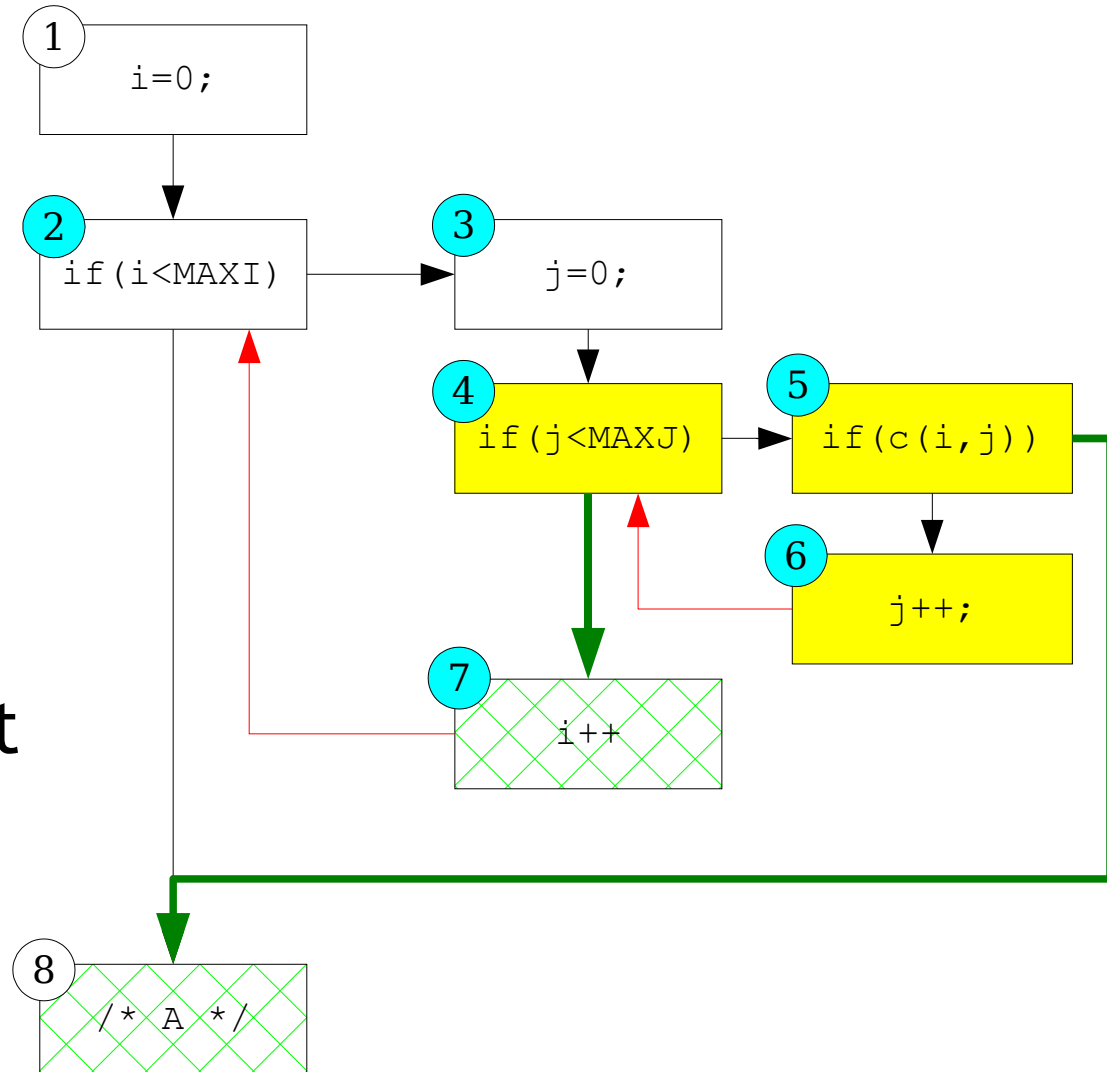
# "Before" the loop

- Always possible
- Inserting a pre-header



# "After" and "around" the loop

- Unique successor:  
unique point after  
(around possible).
- Multiple successors:  
multiple points after  
(around impossible).
- Loops with unique exit  
node allow further  
behaviour prediction.



## ***JP Part 2: Context Exposure (dynamic)***

- Exposing data processed and guiding the execution,
- “Arguments” to the loop,
- Integer range and `Iterators`,
- Arrays and `Collections`.
- (Only loop with unique exit nodes to avoid “break” statements and irregular iterations)

## *Context Exposure*

- For method calls (for example), the context exposed comprises the target, the caller object and the arguments,
- Need similar data for loops to exploit the loop join point potential,
- Otherwise, only able to recognise that there is a loop, but no extra information on what it does.

# *Integer range and Iterators*

- `for (int i = min ; i < max ; i+=1)`
- Need to get *min*, *max* and *stride* for parallelisation.
- `while (iter.hasNext()) { ... iter.next() ... }`
- Need to get Iterator *iter*.
- Passed as “args (min, max, stride)” or “args (iter)”.

# *Arrays and Collection*

- Analogy with Java 5 (Tiger) constructs.
- `for (Object item: collec) { ... }`
- ```
Iterator iter = collec.iterator();
while (iter.hasNext()) {
    Object item = iter.next() ;
    ...
}
```
- Provides extra information about the data processed by the loop.

## *Loop selection*

- In AspectJ, the selection is (ultimately) based on a name pattern, for example on the method name or an argument type,
- Loops haven't got names,
- Selection to be made on argument types and on data processed: integer range and Iterators; and especially arrays and Collections. (`+cflow`, `within` and `withincode`)
- `pointcut` `bytearrayloop(int min, int max, int s, byte[] a) : loop() && args(min, max, s, a);`



## *Implementation using abc*

- `abc`: AspectBench Compiler (full AspectJ compiler),
- *LoopsAJ*: our extension for `abc` that implements a loop pointcut,
- Analysis capabilities of Soot,
- Need to update the graph when weaving,
- Only one “after” point possible,

## *Dealing with exceptions*

- The graph is not necessarily “reducible” (loops may have several entry points),
- The traps for the exceptions do not necessarily match anything in the source code.

## ***Related topics: loop-body join point***

- It would be possible to insert a node similar to the “pre-header”, but for edges from the loop.
- This would comprise the evaluation of the condition within the definition of the “loop-body”.
- What would context could be exposed?

## *Summary*

- Loop join point possible,
- Meaningful thanks to context exposure,
- Problem of loop selection would probably benefit from `pcflow`, `dflow` and even a possible `pdflow`.