

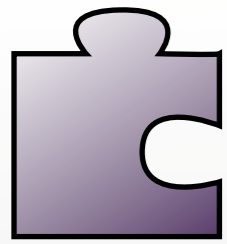
Fitting the Pieces Together: A Machine-Checked Model of Safe Composition

Benjamin Delaware

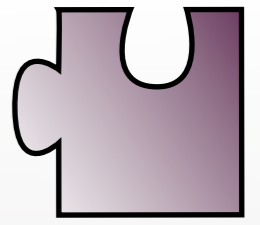
William Cook

Don Batory

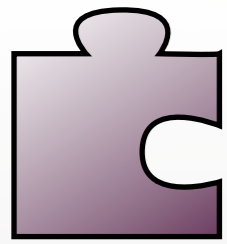
University of Texas at Austin



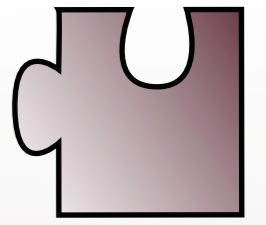
Safe Composition

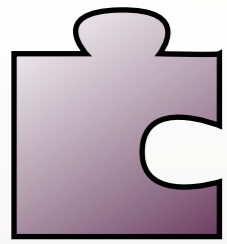


- **Features**
 - Word Processor has formatting, printing, spell check, tables..
 - Cut across traditional modularity boundaries
 - Reify functionality into distinct **feature modules**
- **Software Product Line (SPL)**
 - Multiple products from one code base
 - Product = subset of features
- **Safe Composition**
 - Type check all products
 - Products are exponential in number of features
- **Goal**
 - Sound type system
 - Foundation for efficient implementation

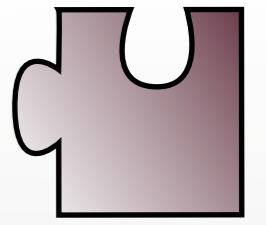


A Feature Example

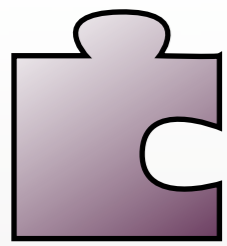




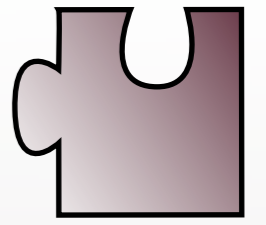
A Feature Example



- Features are sets of class definitions and refinements



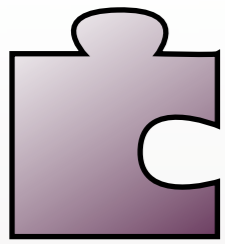
A Feature Example



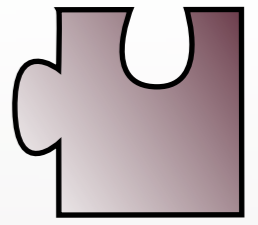
- Features are sets of class definitions and refinements

```
feature Account {  
  class Account extends Object {  
    int balance = 0;  
    void update(int x) {  
      int newBal = balance + x;  
      balance = newBal;  
    }  
  }  
}
```

Account



A Feature Example



- Features are sets of class definitions and refinements

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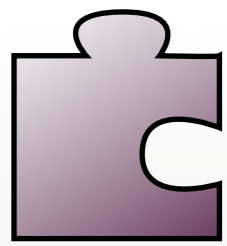
Account

```
feature InvestAccount {  
  refines class Account extends WaMu {  
    int 401kbalance = 0;  
    refines void update (int x) {  
      x = x/2;  
      Super();  
      401kbalance += x;  
    }  
  }  
}
```

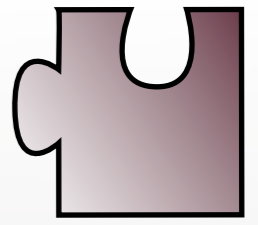
InvestAccount

```
feature RetireAccount {  
  refines class Account extends Lehman {  
    int 401kbalance = 10000;  
    int update (int x) {  
      401kbalance += x;  
    }  
  }  
}
```

RetireAccount



Composing Features



- Features are sets of class definitions and refinements

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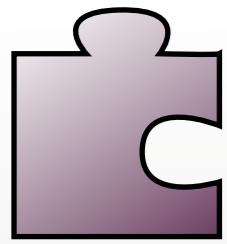
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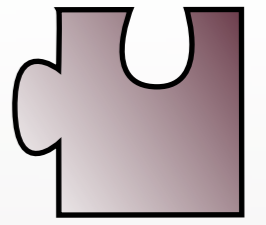
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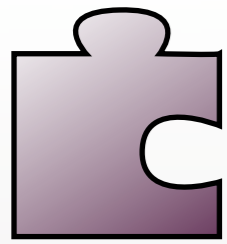
Composing Features



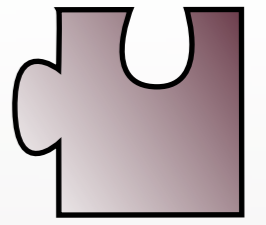
Account

InvestAccount

RetireAccount



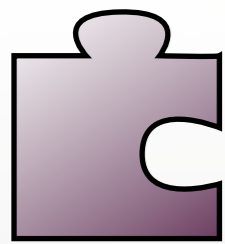
InvestAccount • Investor



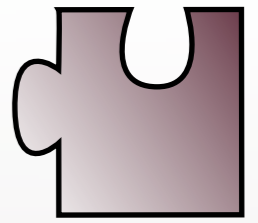
Account

InvestAccount

RetireAccount



InvestAccount • Investor



InvestAccount

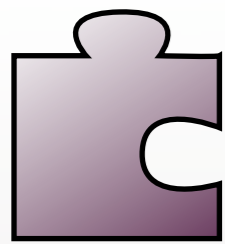
RetireAccount

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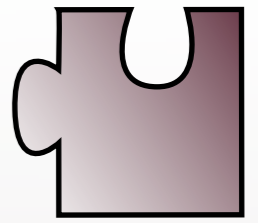
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Account



InvestAccount • Investor



RetireAccount

Account

InvestAccount

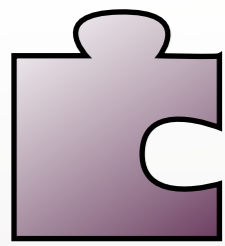
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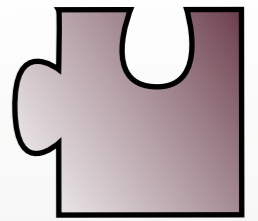
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Account



InvestAccount • Investor



RetireAccount

Account

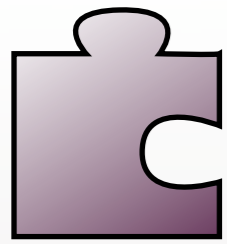
+

InvestAccount

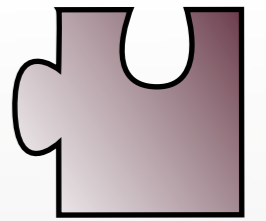
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Account



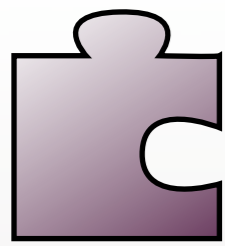
RetireAccount • Investor



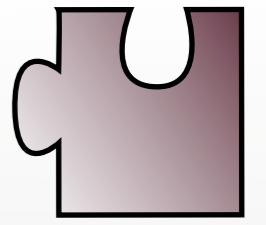
Account

InvestAccount

RetireAccount



RetireAccount • Investor



InvestAccount

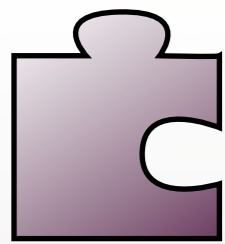
RetireAccount

Account

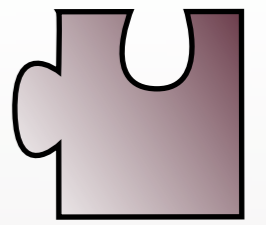
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Account



RetireAccount • Investor



InvestAccount

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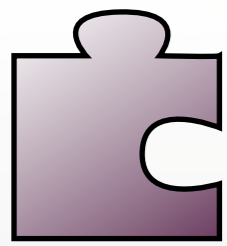
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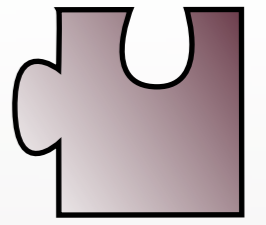
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RetireAccount



RetireAccount • Investor



InvestAccount

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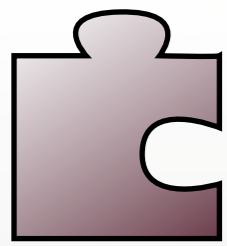
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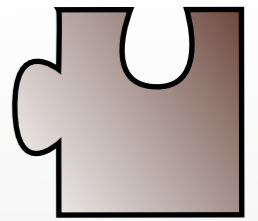
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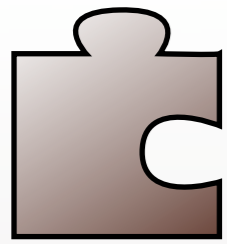
Account



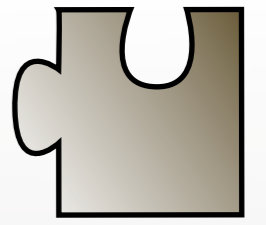
Feature Models



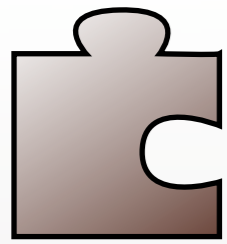
- A SPL has a set of available features:
 $\{\text{Account}, \text{RetireAccount}, \text{InvestAccount}\}$
- Typically feature combinations are constrained
 - A **feature model** represents these constraints
 - Propositional formula is compact representation [Batory05]
 $\text{RetireAccount} \vee \text{InvestAccount}$
 - Product corresponds to truth assignment
- FMs should enforce implementation constraints
 - Safe Composition
 $(\text{RetireAccount} \vee \text{InvestAccount}) \wedge$
 $(\text{RetireAccount} \Rightarrow \text{Account}) \wedge (\text{InvestAccount} \Rightarrow \text{Account})$



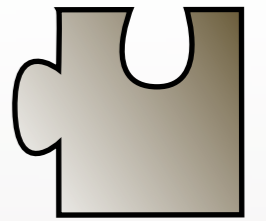
Checking Safe Composition



- Could synthesize entire product line
 - Computationally expensive:

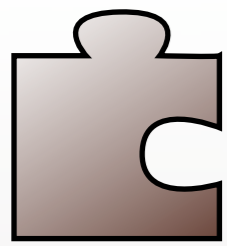


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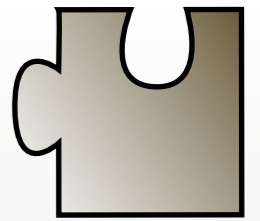


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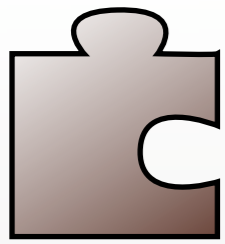


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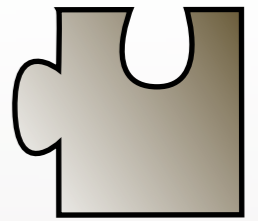


- Could synthesize entire product line
 - Computationally expensive:





Difficulties

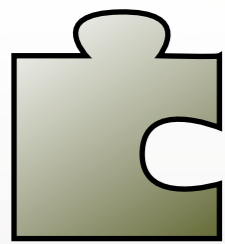


- Combinatorial nature of SPLs problematic:

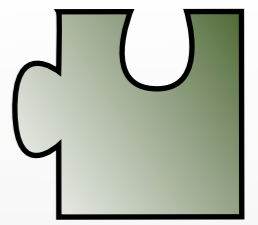
```
feature Payroll {  
  class Employer extends Object {  
    Account Employee1;  
    ...  
    Employee1.401kbalance += 10000;  
    ...  
  }  
}
```

Bailout

- **Bailout** feature needs Account
- Account needs 401kbalance
- Multiple ways to satisfy
 - Introduction
 - Inheritance
- Features are static
 - Surrounding program is not
- Dependencies are resolved by a combination of features
 - These features have their own set of dependencies



Lightweight Feature Java



- Lightweight Java [Strnisa07]
- Minimal imperative subset of Java formalized in Coq
- Lightweight Feature Java
- Lightweight Java extended with features

Feature Table

$FT ::= \{\overline{FD}\}$

Product specification

$PS ::= \overline{F}$

Feature declaration

$FD ::= \mathbf{feature} F \{\overline{cld}; \overline{rcld}\}$

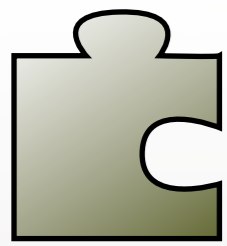
Class refinement

$rcld ::= \mathbf{refines class} dcl \mathbf{extending} cl \{\overline{fd}; \overline{md}; \overline{rmd}\}$

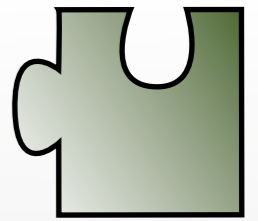
Method Refinement

$rmd ::= \mathbf{refines method} ms \{\overline{s}; \mathbf{Super}(); \overline{s}; \mathbf{return} y\}$

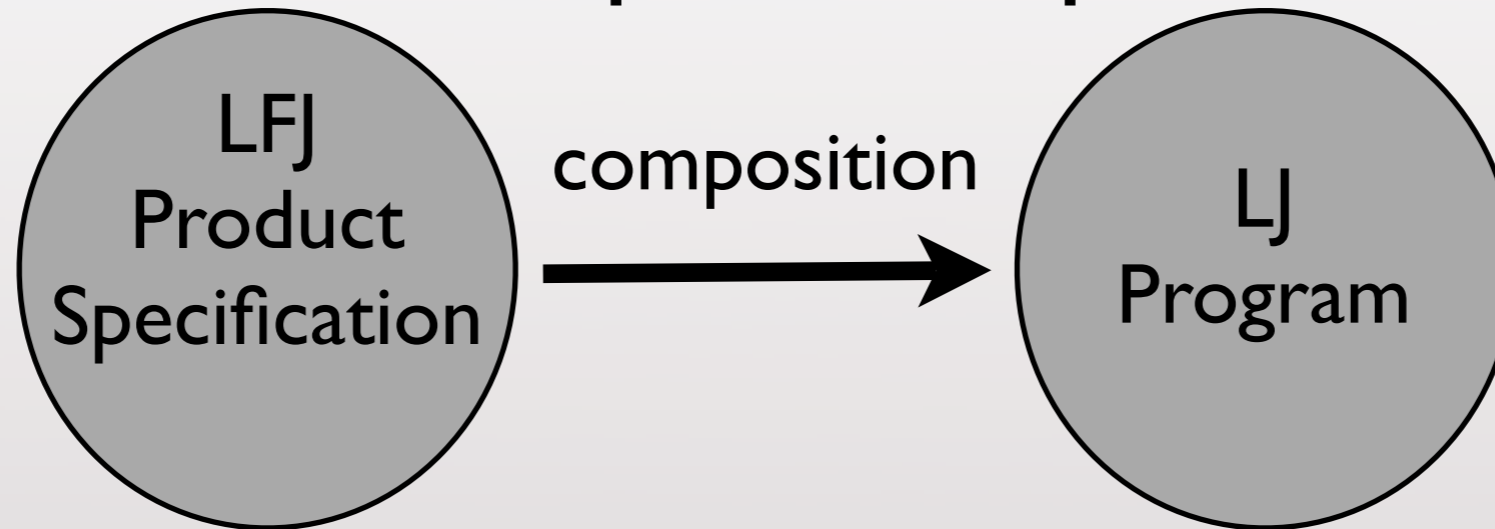
- Formalized in the Coq Proof Assistant



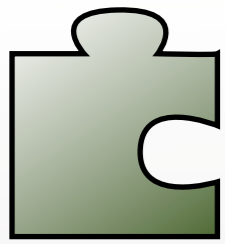
Composition in LFJ



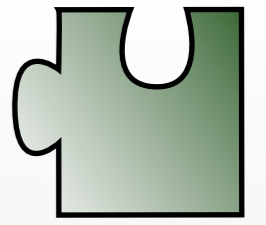
- Programs built from product specifications



- **compose**
 - Refine existing classes
 - Apply method refinement
 - Introduce fields, methods
 - Introduce new classes
- Recursively apply **compose** to specification



LJ Type System

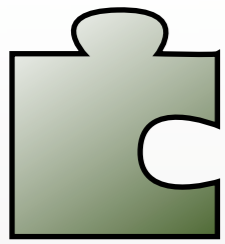

$$\text{distinct}(\overline{var_k^k})$$
$$\frac{}{\text{type}(cl_k) = \tau_k^k}$$
$$\text{type}(cl) = \tau'$$
$$\Gamma = [\overline{var_k^k} \mapsto \tau_k^k][\mathbf{this} \mapsto \tau]$$
$$\Gamma(y) = \tau''$$

(WF-METHOD)

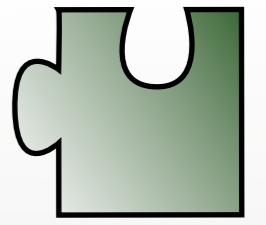
$$\frac{}{\mathbf{P}, \Gamma \vdash s_\ell^\ell}$$
$$\mathbf{P} \vdash \tau'' \prec \tau'$$
$$\frac{}{\mathbf{P} \vdash \text{defined } cl_k^k}$$

$$\mathbf{P} \vdash_\tau cl \text{ meth } (\overline{cl_k^k} \overline{var_k^k}) \{ \overline{s_\ell^\ell} \text{ return } y; \}$$

- Program not available until composition



LJ Type System

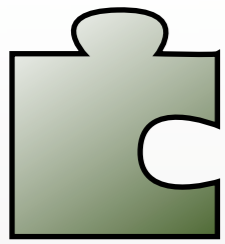

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Internal Checks

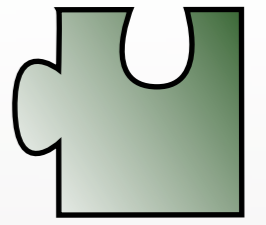
(WF-METHOD)

$$\mathbb{P} \vdash_\tau cl \text{ meth } (\overline{cl_k} \overline{var_k^k}) \{ \overline{s_\ell^\ell} \text{ return } y; \}$$

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LJ Type System



$\text{distinct}(\overline{var_k^k})$

$\frac{}{\text{type}(cl_k) = \tau_k^k}$

$\text{type}(cl) = \tau'$

$\Gamma = [\overline{var_k} \mapsto \tau_k^k][\text{this} \mapsto \tau]$

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(WF-METHOD)

$\frac{}{\mathbf{P}, \Gamma \vdash s_\ell^\ell}$

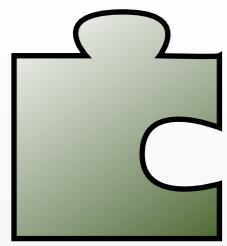
$\mathbf{P} \vdash \tau'' \prec \tau'$

External Checks

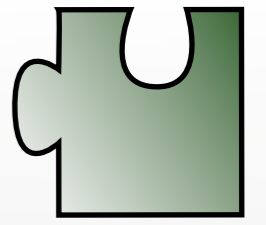
$\frac{}{\mathbf{P} \vdash \text{defined } cl_k^k}$

$\mathbf{P} \vdash_\tau cl \text{ meth } (\overline{cl_k} \overline{var_k^k}) \{ \overline{s_\ell^\ell} \text{ return } y; \}$

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Constraint-Based Typing

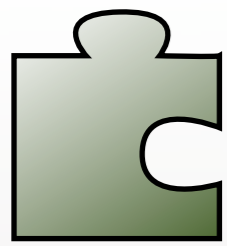


- External premises become constraints

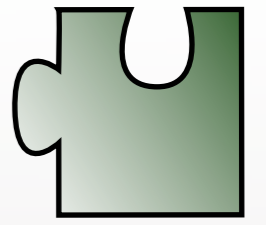
$$\begin{array}{c}
 \text{distinct}(\overline{var_k^k}) \\
 \hline
 \text{type}(cl_k) = \tau_k \\
 \text{type}(cl) = \tau' \\
 \Gamma = [\overline{var_k} \mapsto \tau_k^k][\text{this} \mapsto \tau] \\
 \hline
 \overline{\Gamma \vdash s_\ell \mid \mathcal{C}_\ell} \\
 \Gamma(y) = \tau''
 \end{array}
 \quad (\text{WF-METHOD})$$

$$\vdash_\tau cl \text{ meth } (\overline{cl_k \ var_k^k}) \{ \overline{s_\ell}^\ell \text{ return } y; \} \mid \{ \tau'' \prec \tau', \overline{\text{defined } cl_k^k} \} \cup \cup_\ell \mathcal{C}_\ell$$

- Compositional Constraints
- Uniqueness Constraints
- Structural Constraints



Constraint-Based Typing



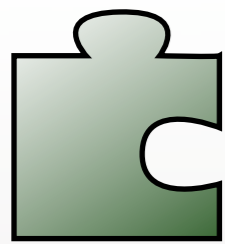
- Two typing phases
- Typing Feature Tables

$$\frac{\vdash \overline{\text{FD}}_k^k \mid \mathbf{WF}_k}{\vdash \{\overline{\text{FD}}_k^k\} \mid \bigcup_k \{\mathbf{In}_{\text{FD}_k} \Rightarrow \mathbf{WF}_k\}}$$

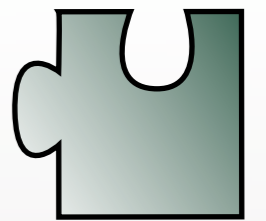
- Well-typed product specification

$$\text{PS} \models \bigcup_k \{\mathbf{In}_{\text{FD}_k} \Rightarrow \mathbf{WF}_k\}$$

- Feature Constraint
- Compositional Constraints
- Uniqueness Constraints
- Structural Constraints



Soundness of LFJ Type System



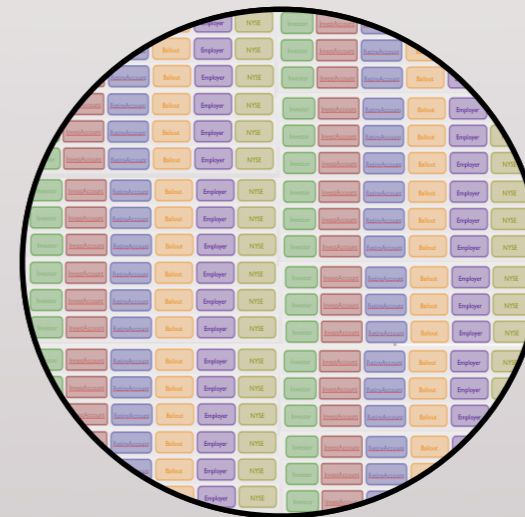
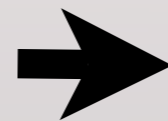
Theorem:

$$\vdash \{FD_k\} \mid \bigcup_k \{\mathbf{In}_{FD_k} \Rightarrow \mathbf{WF}_k\}$$

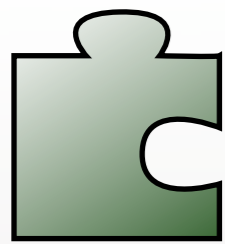
$$PS \models \bigcup_k \{\mathbf{In}_{FD_k} \Rightarrow \mathbf{WF}_k\}$$

$$\vdash_{FJ} \mathbf{compose}(PS)$$

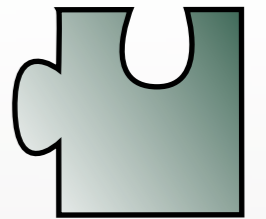
- Space of products



- First premise describes subset of type-safe products
- Second ensures product in this space



Soundness of LFJ Type System



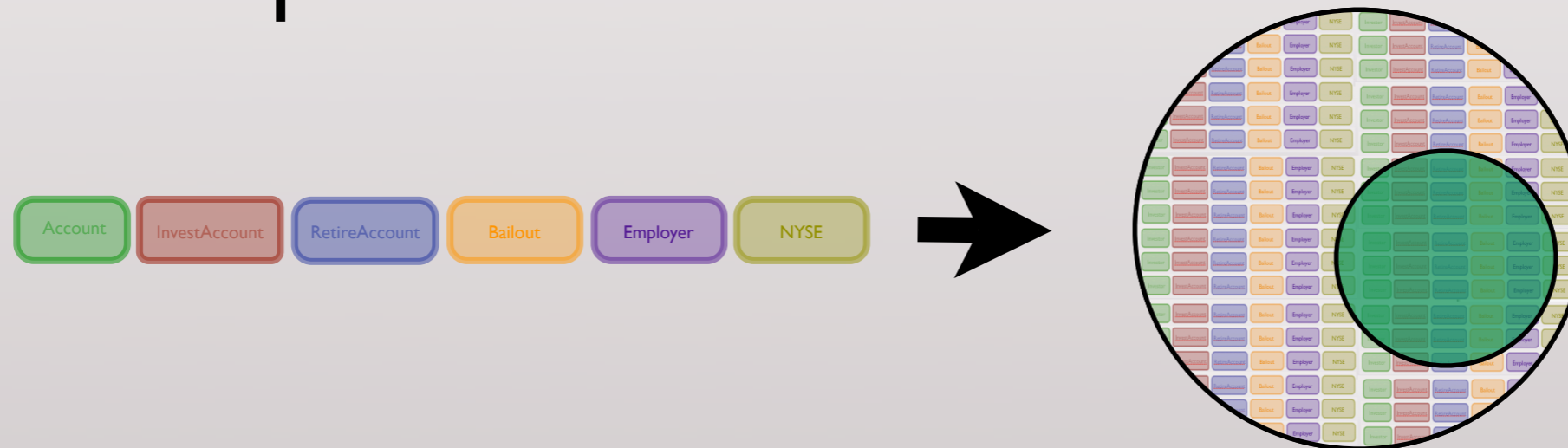
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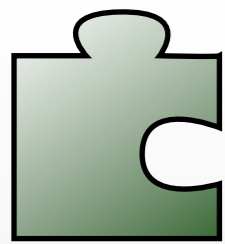
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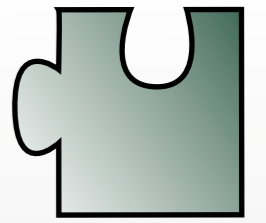
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Soundness of LFJ Type System



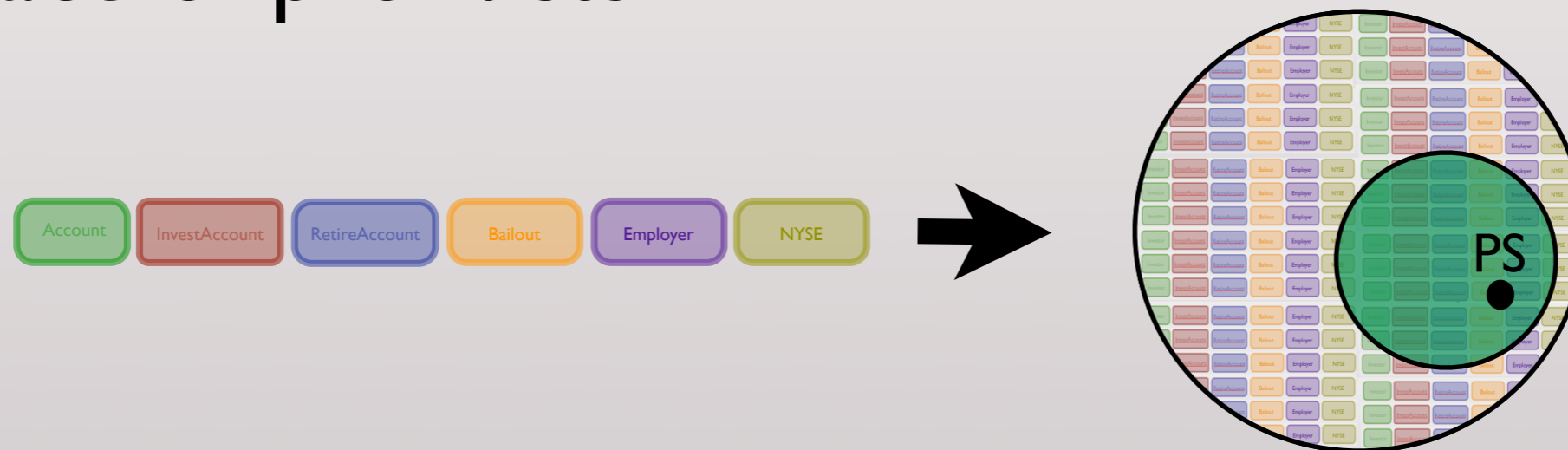
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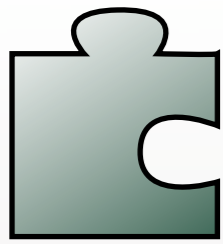
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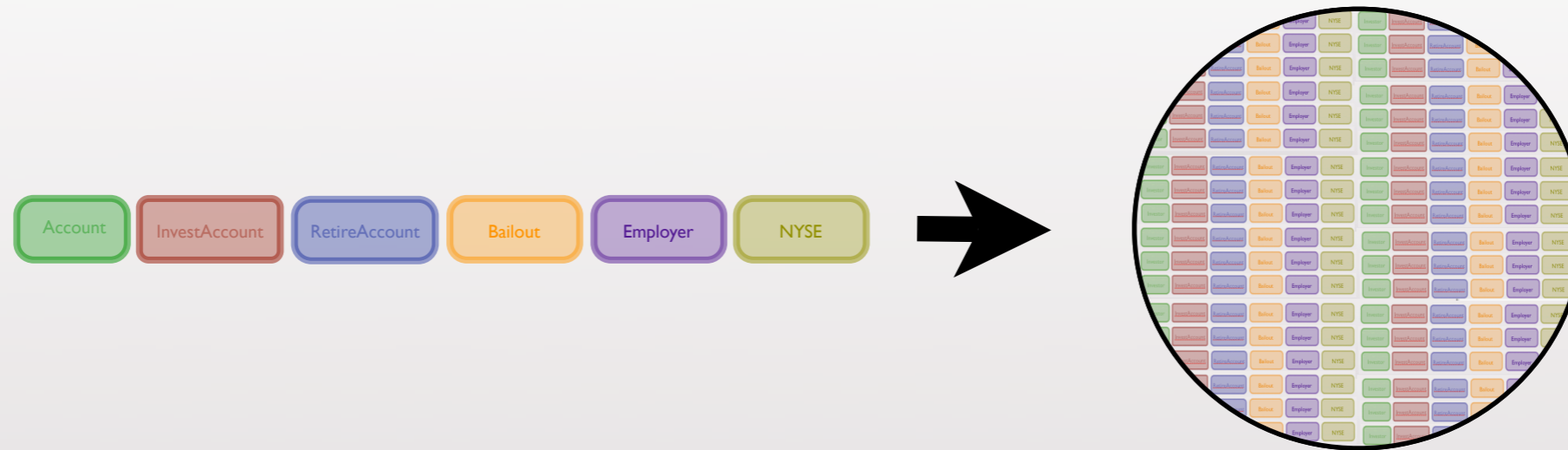
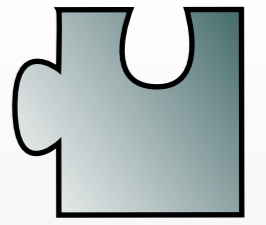
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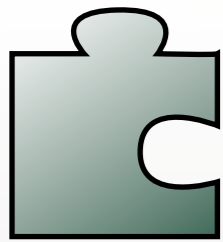


Validating Feature Models

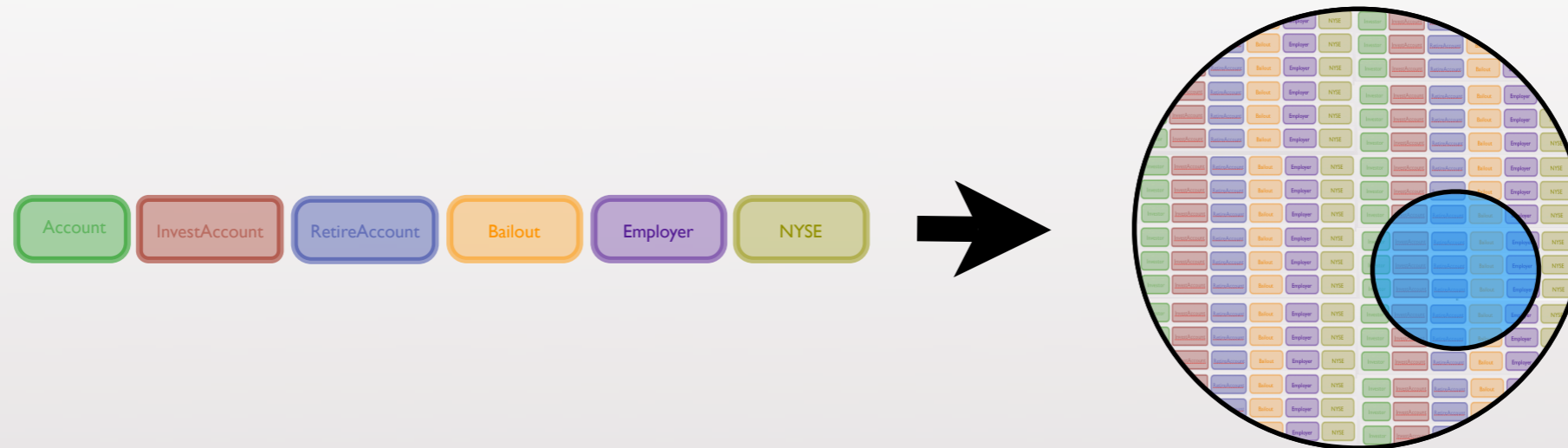
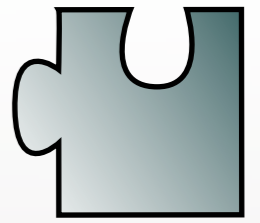


- Feature Models describe desired product space
 - Should be contained in type-safe space
- Recall Feature Models are propositional formulas
 - Describe type-safe space in propositional logic, **WF**_{Safe}
 - Reduction from typing constraints
- Reduce to SAT:

$$\text{FM} \Rightarrow \text{WF}_{\text{Safe}}$$

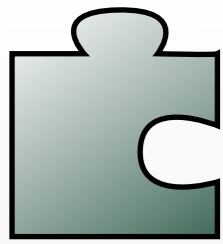


Validating Feature Models

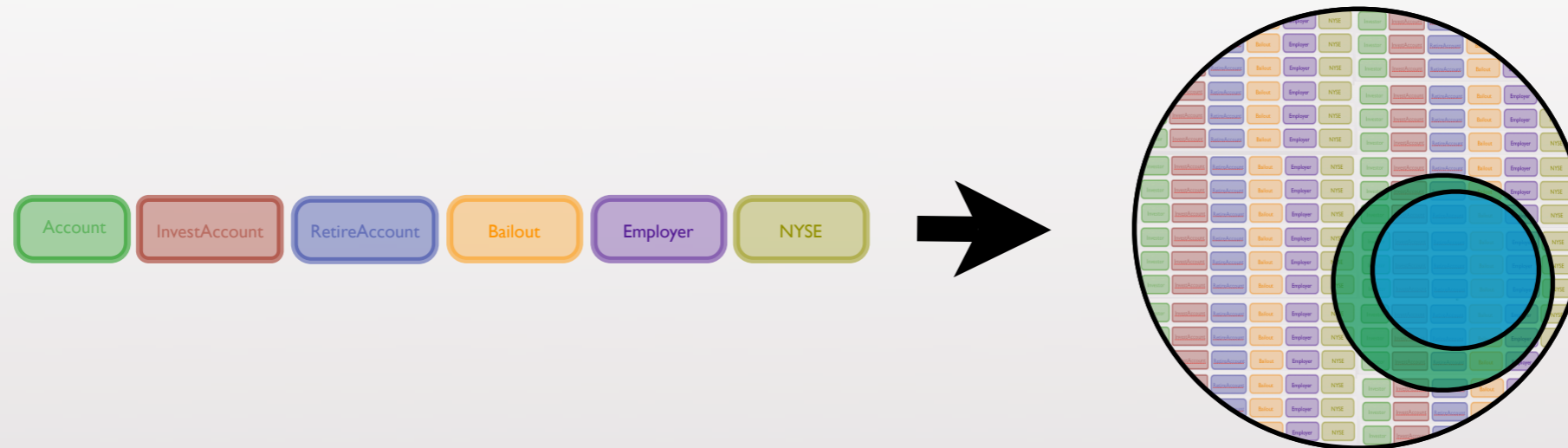
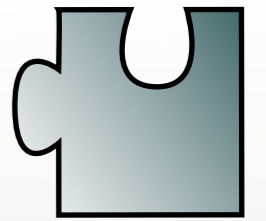


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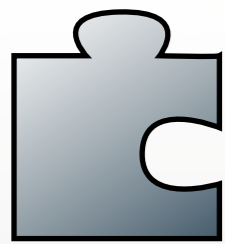


Validating Feature Models

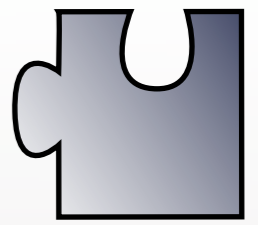


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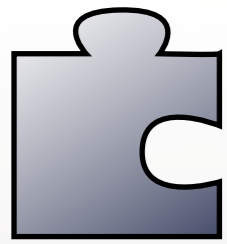
Evaluation



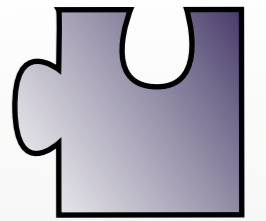
- Checking validity coNP-complete in general
- Our formulas are highly structured

Product Line	# of Features	# of Programs	Code Base Jak/Java LOC	Program Jak/ Java LOC	Typechecking Time
JPL	70	56	34K/48K	22K/35K	<30s

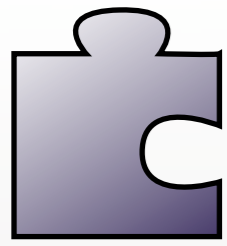
- Previous implementation of approach [Thaker07]
 - Identified errors in existing product lines
- Evidence of erroneous product



Conclusion



- Feature-based Software Product Lines
- Safe Composition
- Lightweight Feature Java
 - Verified in Coq proof assistant
 - Constraints describe program space
- Validating Feature Models
 - Reduce to SAT
 - Efficient evaluation



Questions?

